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Surface Area is roughly at least the volume!

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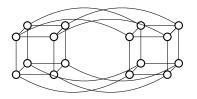
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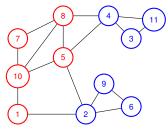
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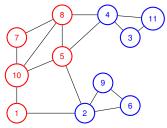
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Restatement: for any cut in the hypercube, the number of cut edges is at least the size of the small side.

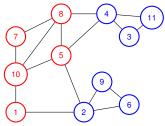


S is red, V - S is blue.



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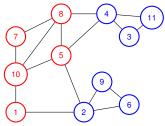
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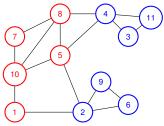
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Number of edges between red and blue. 4.



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Hypercube: any cut that cuts off *x* nodes has $\ge x$ edges.

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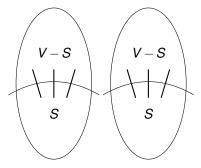
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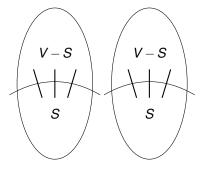
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Case 2: Count inside and across.

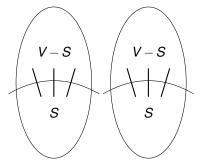


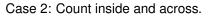
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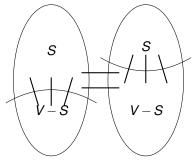
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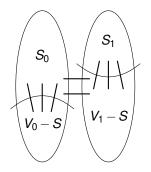
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$$|S_0| \ge |V_0|/2.$$

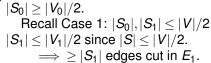


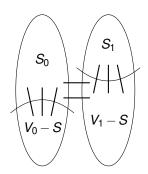
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 S_0 $V_0 - S$ $V_1 - S$ $|S_0| \ge |V_0|/2.$ Recall Case 1: $|S_0|, |S_1| \le |V|/2$ $|S_1| \le |V_1|/2$ since $|S| \le |V|/2.$

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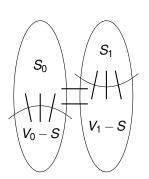
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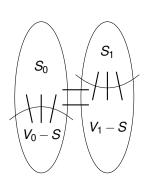
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$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| &\leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| &\leq |V|/2. \\ &\implies &\geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| &\leq |V_0|/2 \end{split}$$

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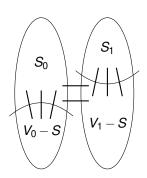
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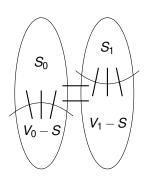


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Edges in E_x connect corresponding nodes.

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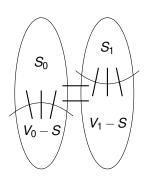


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Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

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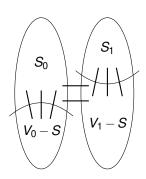


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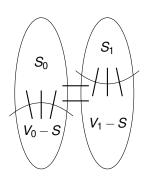
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Total edges cut:

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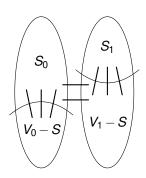
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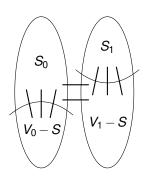
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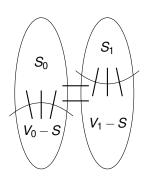
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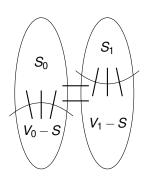
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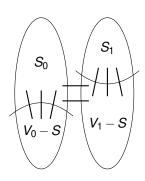
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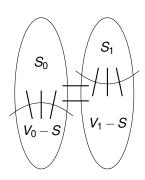
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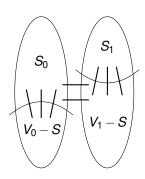
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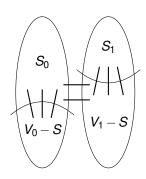
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Also, case 3 where $|S_1| \ge |V|/2$ is symmetric.

Hypercube proof: poll

Hypercube has large cuts proof uses these ideas:

- (A) If cuts are same size on two sides it works by induction.
- (B) Uses the fact that it is planar.
- (C) Recursive definition of hypercube.
- (D) If different size, can count edges between to subcubes.
- (E) Applies Euler's formula.

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(A),(D), and (E).

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Central object of study.

Modular Arithmetic.

Applications: cryptography, error correction.

Theorem: If d|x and d|y, then d|(y-x).

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Proof: x = ad, y = bd, $(x - y) = (ad - bd) = d(a - b) \implies d|(x - y).$

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Theorem: Every number $n \ge 2$ can be represented as a product of primes.

Proof: Either prime, or $n = a \times b$, and use strong induction. (Uniqueness? Later.)

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
- (B) Euler's formula.
- (C) The definition of a prime number.
- (D) Euclid's Lemma.

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
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- (C) The definition of a prime number.
- (D) Euclid's Lemma.
- (A) and (C)

Next Up.

Modular Arithmetic.

If it is 1:00 now.

If it is 1:00 now. What time is it in 2 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system.

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What time is it in 100 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

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What time is it in 100 hours? 101:00!

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What time is it in 100 hours? 101:00! or 5:00.

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5 is the same as 101 for a 12 hour clock system.

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Clock time equivalent up to addition of any integer multiple of 12.

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Custom is only to use the representative in $\{12, 1, \dots, 11\}$

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Custom is only to use the representative in $\{12, 1, ..., 11\}$ (Almost remainder, except for 12 and 0 are equivalent.)

This is Thursday is February 11, 2021.

This is Thursday is February 11, 2021. What day is it a year from then?

This is Thursday is February 11, 2021. What day is it a year from then? on February 11, 2022?

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5 days from then.

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5 days from then. day 9

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5 days from then. day 9 or day 2

This is Thursday is February 11, 2021. What day is it a year from then? on February 11, 2022? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

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two days are equivalent up to addition/subtraction of multiple of 7.

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11 days from then is day 1

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What day is it a year from then? Next year is not a leap year.

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What day is it a year from then?

Next year is not a leap year. So 365 days from then.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7. divide and get remainder.

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Today: day 4.

5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7. divide and get remainder. 369/7

This is Thursday is February 11, 2021. What day is it a year from then? on February 11, 2022? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

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Next year is not a leap year. So 365 days from then.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5.

This is Thursday is February 11, 2021. What day is it a year from then? on February 11, 2022? Number days.

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Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

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or February 11, 2022 is a Friday.

This is Thursday is February 11, 2021. What day is it a year from then? on February 11, 2022? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

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or February 11, 2022 is a Friday.

80 years?

80 years? 20 leap years.

80 years? 20 leap years. 366×20 days

80 years? 20 leap years. 366×20 days 60 regular years.

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80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$.

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Hmm.

What is remainder of 366 when dividing by 7?

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

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Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7?

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Years and years...
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Years and years...
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Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60$

```
Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$

```
Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7?

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
60 regular years. 365 \times 60 days
Today is day 4.
It is day 4 + 366 \times 20 + 365 \times 60. Equivalent to?
```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7
```

```
Years and years...
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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
```

```
Years and years...
```

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80 years? 20 leap years. 366 \times 20 days
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Today is day 4.
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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 4.
Get Day: 4+2 \times 20+1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
Or February 11, 2101 is Saturday!
```

```
Years and years...
```

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60 regular years. 365 \times 60 days
Today is day 4.
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```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
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Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
Or February 11, 2101 is Saturday!
```

Further Simplify Calculation:

```
Years and years...
```

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What is remainder of 365 when dividing by 7? 1
Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
```

Or February 11, 2101 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

```
Years and years...
```

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Remainder when dividing by 7? $104 = 14 \times 7 + 6$. Or February 11, 2101 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $2+2\times 6+1\times 4=18$.

```
Years and years...
```

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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
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Or February 11, 2101 is Saturday!
```

```
Further Simplify Calculation:
```

```
20 has remainder 6 when divided by 7.
60 has remainder 4 when divided by 7.
Get Day: 2+2\times 6+1\times 4=18.
Or Day 6.
```

```
Years and years...
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Or February 11, 2101 is Saturday!

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Further Simplify Calculation:
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20 has remainder 6 when divided by 7.

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Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 6. February 11, 2101 is Saturday.

"Reduce" at any time in calculation!

x is congruent to y modulo m or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by m.

x is congruent to y modulo m or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by m.

... or x and y have the same remainder w.r.t. m.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Mod 7 equivalence classes:

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*. Mod 7 equivalence classes:

woo / equivalence classes

 $\{\ldots, -7, 0, 7, 14, \ldots\}$

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Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\}$

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Mod 7 equivalence classes:

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Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

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 $\{\ldots, -7, 0, 7, 14, \ldots\}$ $\{\ldots, -6, 1, 8, 15, \ldots\}$...

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or " $a \equiv c \pmod{m}$ and $b \equiv d \pmod{m}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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$$a \equiv c \pmod{m}$$
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$$a \equiv c \pmod{m}$$
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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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$$a \equiv c \pmod{m}$$
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 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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$$a \equiv c \pmod{m}$$
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 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore,

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a + b = c + d + (k + j)m and since k + j is integer.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m and since k+j is integer. $\implies a+b \equiv c+d \pmod{m}$.

Can calculate with representative in $\{0, \ldots, m-1\}$.

x (mod m) or mod(x,m)

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

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 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

 $mod(x,m) = x - \lfloor \frac{x}{m} \rfloor m$

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by $m in \{0, ..., m-1\}$.

mod $(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

 $mod (29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12$

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = 4

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29, 12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = ¥ = 5

x (mod m) or mod (x,m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system.

x (mod *m*) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29, 12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = \cancel{x} = 5 Work in this system. $a \equiv b \pmod{m}$.

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system. $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system. 2 = b (mod m)

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system. $a = b \pmod{m}$

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Inverses and Factors.

Division: multiply by multiplicative inverse.

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"Common factor of 4"

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"Common factor of 4" \implies 8k - 12l is a multiple of four for any l and k \implies 8k \neq 1 (mod 12) for any k.

Poll

Mark true statements.

(A) Mutliplicative inverse of 2 mod 5 is 3 mod 5.

- (B) The multiplicative inverse of $((n-1) \pmod{n} = ((n-1) \pmod{n})$.
- (C) Multiplicative inverse of 2 mod 5 is 0.5.
- (D) Multiplicative inverse of $4 = -1 \mod 5$.
- (E) (-1)x(-1) = 1. Woohoo.

(F) Multiplicative inverse of 4 mod 5 is 4 mod 5.

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(F) Multiplicative inverse of 4 mod 5 is 4 mod 5.

(C) is false. 0.5 has no meaning in arithmetic modulo 5.

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Each of *m* numbers in *S* correspond to one of *m* equivalence classes modulo *m*.

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Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

 \implies $(a-b) \ge m$. But $a, b \in \{0, ..., m-1\}$. Contradiction.

Proof review. Consequence.

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 What is x? Multiply both sides by 5.
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Very different for elements with inverses.

Proof Review 2: Bijections.

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Not a bijection.

Poll

Which is bijection?

(A) f(x) = x for domain and range being \mathbb{R} (B) f(x) = axmod(n) for $x \in \{0, ..., n-1\}$ and gcd(a, n) = 2(C) f(x) = axmodn for $x \in \{0, ..., n-1\}$ and gcd(a, n) = 1

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Find gcd (x, m).

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Find gcd (x, m). Greater than 1?

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Algorithm:

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Finding inverses.

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Next up.

Next up.

Next up. Euclid's Algorithm.

Next up.

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Divisibility...

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Modular Arithmetic Lecture in a minute.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

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For $a \equiv b \pmod{N}$, and $c \equiv d \pmod{N}$, $ac = bd \pmod{N}$ and $a+b=c+d \pmod{N}$.

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Know if there is an inverse, but how do we find it?

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