

Today.

Polynomials.

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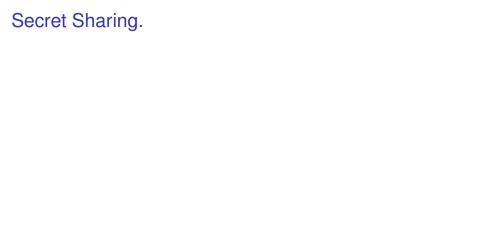
Secret Sharing.

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Polynomials.

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Correcting for loss or even corruption.



Share secret among \boldsymbol{n} people.

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Secrecy: Any k-1 knows nothing.

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The idea of the day.

Two points make a line. Lots of lines go through one point.

A polynomial

$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0.$$

is specified by **coefficients** $a_d, \dots a_0$.

¹A field is a set of elements with addition and multiplication operations, with inverses. $GF(p) = (\{0, ..., p-1\}, + \pmod{p}, * \pmod{p}).$

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Polynomials over reals: $a_1, \ldots, a_d \in \Re$, use $x \in \Re$.

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Polynomials P(x) with arithmetic modulo p: ¹ $a_i \in \{0, ..., p-1\}$ and

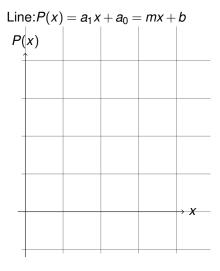
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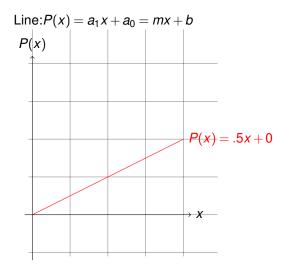
for $x \in \{0, ..., p-1\}$.

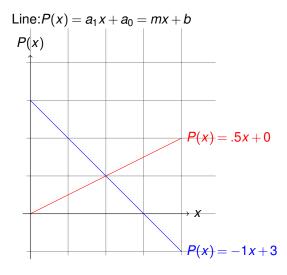
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Parabola: $P(x) = a_2x^2 + a_1x + a_0$

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Parabola: $P(x) = a_2x^2 + a_1x + a_0 = ax^2 + bx + c$

Line:
$$P(x) = a_1 x + a_0 = mx + b$$

$$P(x) = 0.5x^2 - x + 0.1$$

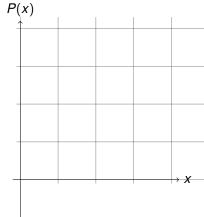
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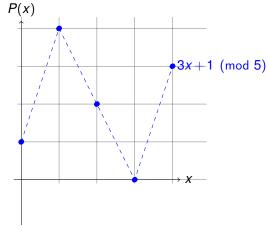
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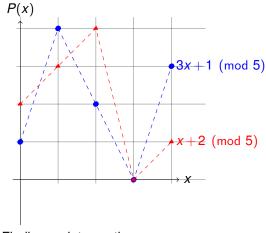
$$P(x) = 0.5x^2 - x + 0.1$$

$$P(x) = -.3x^2 + 1x + .1$$

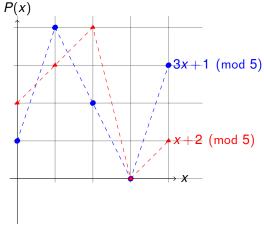
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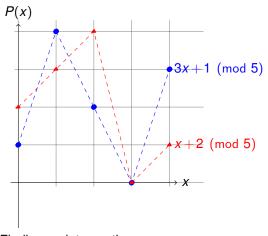




Finding an intersection.
$$x+2 \equiv 3x+1 \pmod{5}$$
 $\implies 2x \equiv 1 \pmod{5}$



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$$x+2\equiv 3x+1\pmod{5}$$
 $\implies 2x\equiv 1\pmod{5}$ $\implies x\equiv 3\pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Good when modulus is prime!!

Fact: Exactly 1 degree $\leq d$ polynomial contains d+1 points. ²

²Points with different *x* values.

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Poll.

Two points determine a line. What facts below tell you this?

Say points are $(x_1, y_1), (x_2, y_2)$ **.**

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- (A) Line is y = mx + b.
- (B) Plug in a point gives an equation: $y_1 = mx_1 + b$
- (C) The unknowns are *m* and *b*.
- (D) If equations have unique solution, done.

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All true.

Flow Poll.

Why solution? Why unique?

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- (A) Solution cuz: $m = (y_2 y_1)/(x_2 x_1), b = y_1 m(x_1)$
- (B) Unique cuz, only one line goes through two points.
- (C) Try: $(m'x + b') (mx + b) = (m' m)x + (b b) = ax + b \neq 0$.
- (D) Either $ax_1 + c \neq 0$ or $ax_2 + c \neq 0$.
- (E) Contradiction.

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Flow poll. (All true. (B) is not a proof, it is restatement.)

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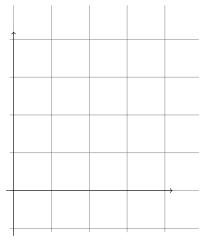
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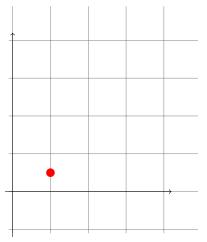
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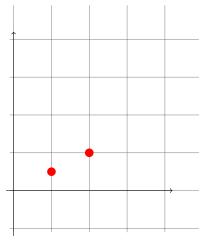
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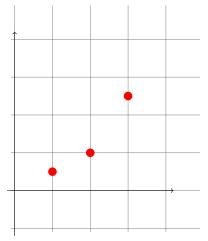
Fact: Exactly 1 degree $\leq d$ polynomial contains d+1 points. ³



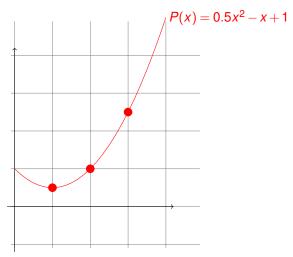
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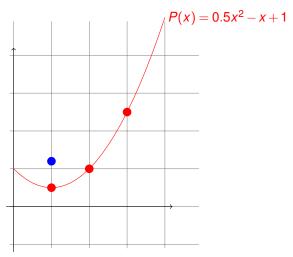
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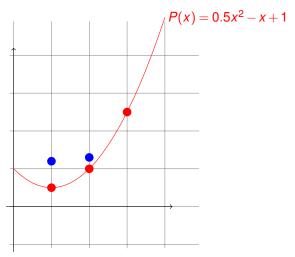
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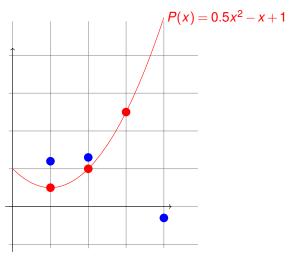
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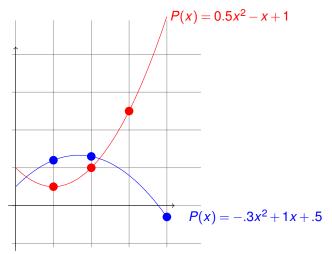
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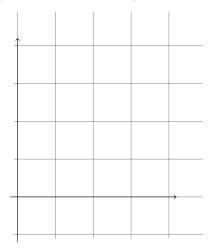


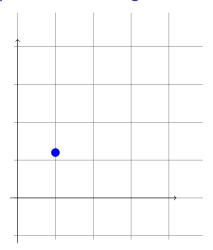
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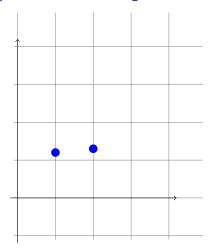


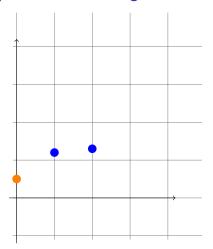
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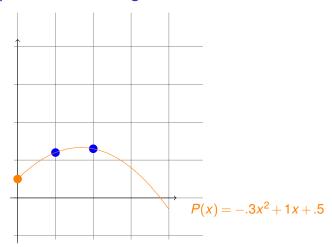
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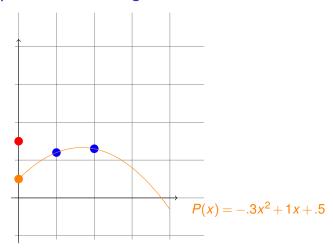


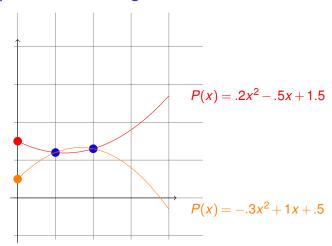


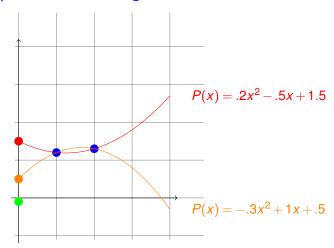


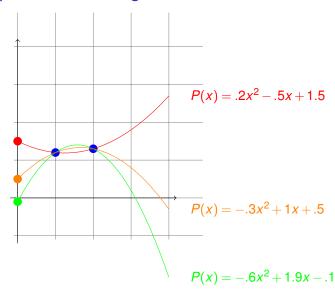


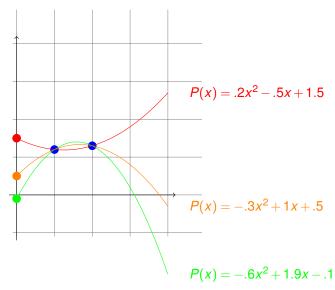












Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Knowing k pts \implies only one P(x) \implies evaluate P(0).

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For a line, $a_1x + a_0 = mx + b$ contains points (1,3) and (2,4).

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And the line is...

$$x+2 \mod 5$$
.

For a quadratic polynomial, $a_2x^2 + a_1x + a_0$ hits (1,2); (2,4); (3,0).

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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

In general..

Given points: (x_1, y_1) ; $(x_2, y_2) \cdots (x_k, y_k)$.

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Multiplicative inverses due to gcd(x,p) = 1, for all $x \in \{1,...,p-1\}$

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See the idea?

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Given d+1 points, use Δ_i functions to go through points? $(x_1, y_1), \ldots, (x_{d+1}, y_{d+1}).$

Will $y_1 \Delta_1(x)$ contain (x_1, y_1) ?

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See the idea? Function that contains all points?

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Mark what's true.

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Proof:

Assume two different polynomials Q(x) and P(x) hit the points.

Uniqueness Fact. At most one degree d polynomial hits d+1 points.

Roots fact: Any nontrivial degree *d* polynomial has at most *d* roots.

Non-zero line (degree 1 polynomial) can intersect y = 0 at only one x.

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That is, $P(x)=(x-a)Q(x)+r$

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Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.

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(Almost) the same as what is missing: one P(i).



Runtime.

Runtime: polynomial in k, n, and $\log p$.

- 1. Evaluate degree k-1 polynomial n times using $\log p$ -bit numbers.
- 2. Reconstruct secret by solving system of *k* equations using $\log p$ -bit arithmetic.

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Infinite number for reals, rationals, complex numbers!

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Solution: lagrange interpolation.

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Apply: P(x), Q(x) degree d.

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Secret Sharing:

k points on degree k-1 polynomial is great! Can hand out n points on polynomial as shares.